

# Pattern avoidance in words

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## Definition

A **word** on a set  $\Sigma$  is a sequence of elements from  $\Sigma$ .

For this talk:  $\Sigma = \{0, 1\}$  (the “alphabet”)

Example:  $w = 0110$

Pattern containment:

## Example

0110 **contains** 11 (as a contiguous subword).

0110 **avoids** 00.

How many length- $n$  words avoid 01?

length 0:	$\epsilon$	1
length 1:	0, 1	2
length 2:	00, 10, 11	3
length 3:	000, 100, 110, 111	4
	$\vdots$	$\vdots$
length $n$ :		$n + 1$

How many length- $n$  words avoid 10?

length 3:	000, 001, 011, 111	4
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Also  $n + 1$ .

There is a **bijection** from words avoiding 01 to words avoiding 10.  
(Actually two bijections!)

How many length- $n$  words avoid 00?

length 0:	$\epsilon$	1
length 1:	0, 1	2
length 2:	01, 10, 11	3
length 3:	010, 011, 101, 110, 111	5
length 4:	0101, 0110, 0111, 1010, 1011, 1101, 1110, 1111	8
$\vdots$		$\vdots$
length $n$ :		$F(n + 2)$

How many length- $n$  words avoid 11? Also  $F(n + 2)$ .

## Definition

Two patterns are **avoiding-equivalent** if they are avoided by the same number of length- $n$  words for all  $n \geq 0$ .

Length-2 patterns come in two equivalence classes:  $\{01, 10\}$  and  $\{00, 11\}$ .

What are the equivalence classes of length-3 patterns?

pattern $p$	number of length- $n$ words avoiding $p$ for $n = 0, 1, \dots$
000	1, 2, 4, 7, 13, 24, 44, 81, 149, 274, 504, ...
001	1, 2, 4, 7, 12, 20, 33, 54, 88, 143, 232, ...
010	1, 2, 4, 7, 12, 21, 37, 65, 114, 200, 351, ...
011	1, 2, 4, 7, 12, 20, 33, 54, 88, 143, 232, ...
100	1, 2, 4, 7, 12, 20, 33, 54, 88, 143, 232, ...
101	1, 2, 4, 7, 12, 21, 37, 65, 114, 200, 351, ...
110	1, 2, 4, 7, 12, 20, 33, 54, 88, 143, 232, ...
111	1, 2, 4, 7, 13, 24, 44, 81, 149, 274, 504, ...

Equivalence classes:  $\{000, 111\}$ ,  $\{010, 101\}$ ,  $\{001, 011, 100, 110\}$

No surprises.

What are the equivalence classes of length-4 patterns?

equivalence class	self-overlap lengths
$\{0000, 1111\}$	1, 2, 3, 4
$\{0101, 1010\}$	2, 4
$\{0010, 0100, 0110, 1001, 1011, 1101\}$	1, 4
$\{0001, 0011, 0111, 1000, 1100, 1110\}$	4

Surprise!

Where does **0110** overlap itself?

<b>0</b> 110	<b>0</b> 110	<b>0</b> 110	0110
011 <b>0</b>	011 <b>0</b>	011 <b>0</b>	011 <b>0</b>

**Theorem (Solovyov 1966, Guibas 1979)**

*If two patterns have the same set of self-overlap lengths, then they are avoiding-equivalent.*

Proof: Generating series.

Is there a bijective proof?

## Question

If two patterns  $p, q$  are avoiding-equivalent, is there a *natural* bijection from the length- $n$  words avoiding  $p$  to the length- $n$  words avoiding  $q$ ?

length-4 words avoiding  $p = 011$ :

0000  
0001  $\mapsto$  0111  
0010  $\mapsto$  0110  
0100  
0101  
1000  
1001  $\mapsto$  1011  
1010  
1100  
1101  
1110  
1111

length-4 words avoiding  $q = 001$ :

0000  
0100  
0101  
0110  
0111  
1000  
1010  
1011  
1100  
1101  
1110  
1111

Idea: Replace all instances of  $q$  with  $p$ .

0001  $\mapsto$  0011  $\mapsto$  0111

Let  $\phi_L$  be the map that iteratively replaces the leftmost  $q$  with  $p$ .

Let  $A_n(p)$  denote the set of length- $n$  words that avoid  $p$ .

### Lemma

*If  $p, q$  are equal-length patterns and  $w \in A_n(p)$ , then  $\phi_L(w) \in A_n(q)$ .*

A **border** of  $p$  is a subword that is both a prefix and a suffix of  $p$ .

### Theorem

*Let  $p, q$  be patterns such that every nontrivial border of  $p$  is also a border of  $q$ . Then  $\phi_L$  is a bijection from  $A_n(p)$  to  $A_n(q)$ .*

$p = 011$  and  $q = 001$  have no nontrivial borders. ✓

### Non-example

$p = 0100$  and  $q = 1011$  both have length-1 borders, but  $0 \neq 1$ .

$\phi_L(1011110) = 0100110$ .

$\phi_L(0101011) = 0100100 = \phi_L(1011100)$ .      Not a bijection!