

# Morphic words governing the boundaries of cellular automata

Charles Brummitt<sup>1</sup>   Eric Rowland<sup>2</sup>

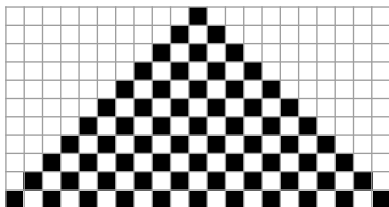
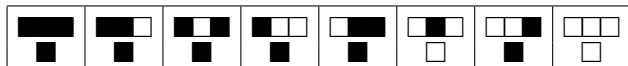
<sup>1</sup>University of California, Davis

<sup>2</sup>LaCIM, Université du Québec à Montréal

May 23, 2012

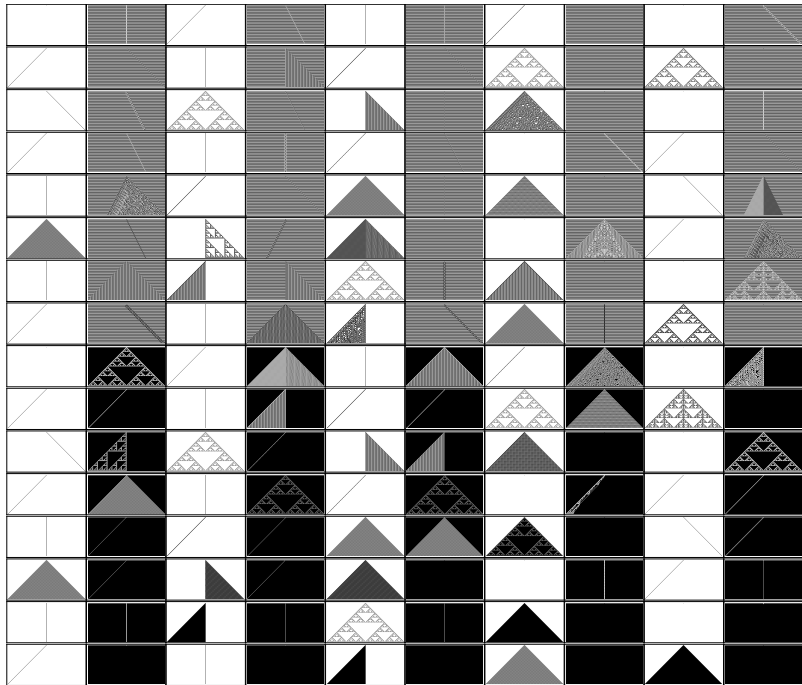
# One-dimensional cellular automata

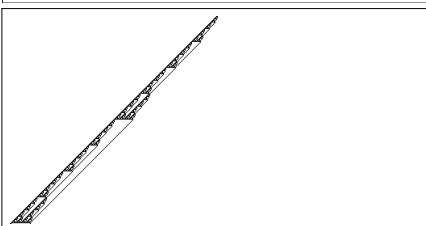
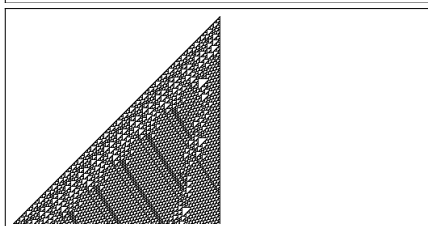
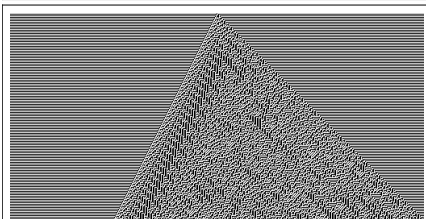
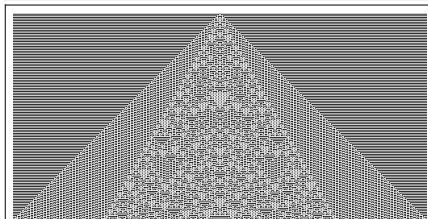
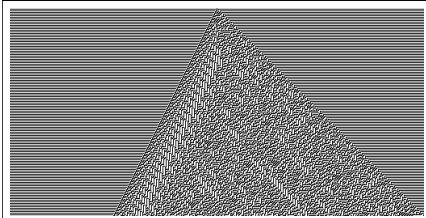
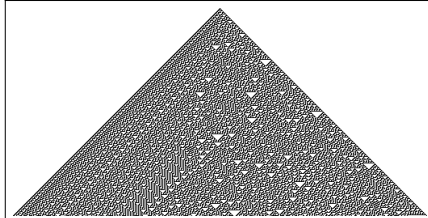
- alphabet  $\Sigma$  of size  $k$  (for example  $\{0, 1, \dots, k - 1\}$ )
- function  $i : \mathbb{Z} \rightarrow \Sigma$  (the initial condition)
- function  $f : \Sigma^d \rightarrow \Sigma$  (the update rule)



Naming scheme:  $11111010_2 = 250$ .

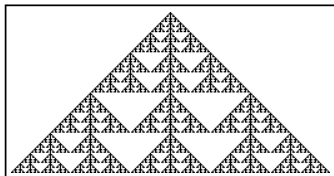
Wolfram: Look at all  $k^{k^d}$   $k$ -color rules depending on  $d$  cells.





## Definition

$\ell(n)$  = width of region on row  $n$  that differs from the background



For example,  $\ell(n) = 2n + 1$ .

Upper bound:  $\ell(n) \leq (d - 1)n + c$ .

For many automata,  $\ell(n)$  grows linearly.

Growth of the form

$$\ell(n) = \begin{cases} an + c_0 & \text{if } n \equiv 0 \pmod{m} \\ an + c_1 & \text{if } n \equiv 1 \pmod{m} \\ \vdots & \vdots \\ an + c_{m-1} & \text{if } n \equiv m-1 \pmod{m} \end{cases}$$

for  $n \geq n_{\min}$  (for some  $m, n_{\min} \in \mathbb{Z}$  and  $a, c_i \in \mathbb{Q}$ )

can be detected by determining whether the difference sequence

$$\ell(n+1) - \ell(n)$$

is eventually periodic.

# Automated–manual search

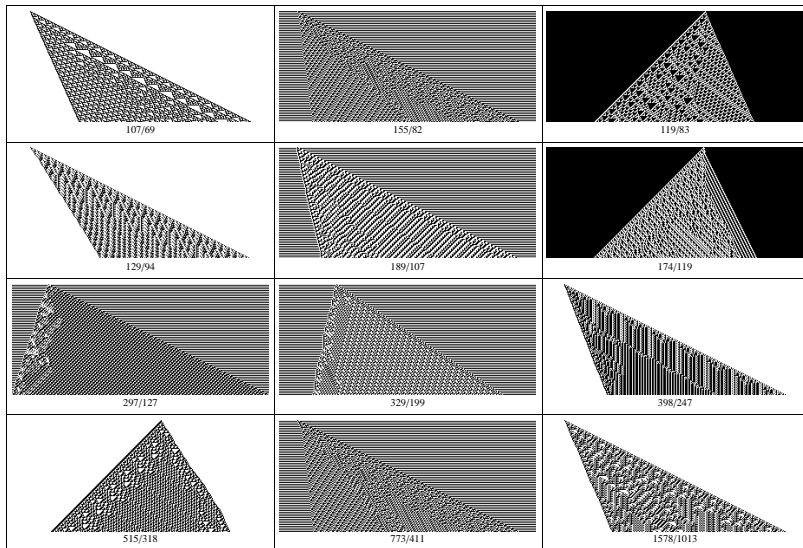
For  $k = 2$  and  $d \leq 3$ , the only growth rates that occur are  $0, 1, 3/2, 2$ .

What happens for larger  $d$ ?

There are  $2^{2^4} = 65536$  2-color rules depending on  $d = 4$  cells.

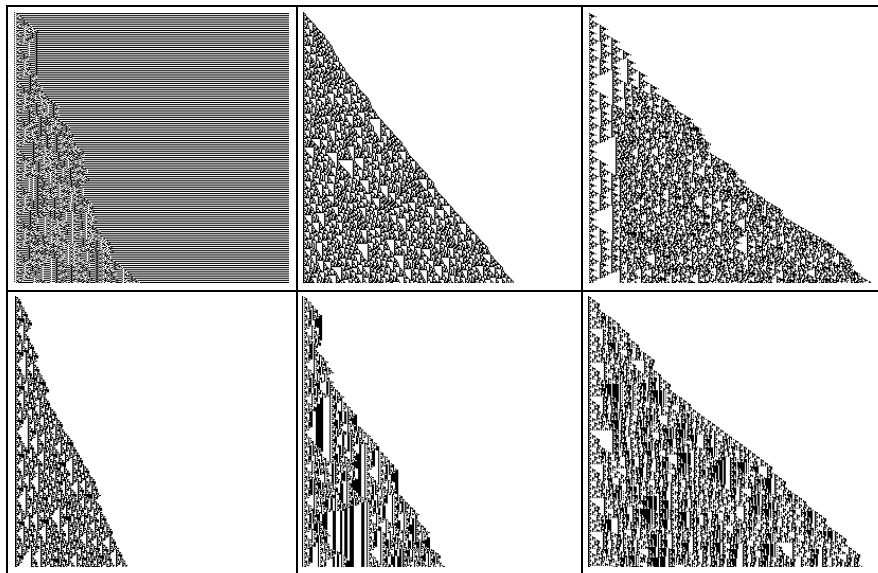
- Consider each rule begun from  $\dots \square\square\square\square \blacksquare\square\square\square \dots$   
and  $\dots \blacksquare\blacksquare\blacksquare\blacksquare \square \blacksquare\blacksquare\blacksquare\blacksquare \dots$ . There are 25088 equivalence classes.
- Filter out automata with sequences  $\ell(n+1) - \ell(n)$  that seem to be eventually periodic.
- Examine the remaining automata manually for reducibility.
- For the remaining automata, fit curves to  $\ell(n)$ .

# The most complex growth rates for $d = 4$

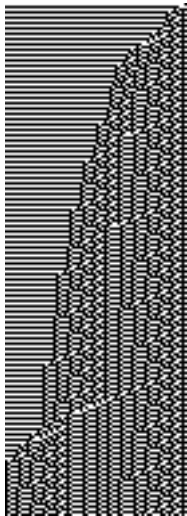




# Chaotic boundaries

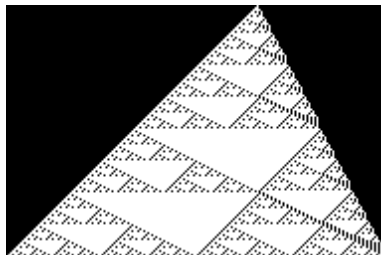


# A misleading automaton



Around step 524500,  
growth increases dramatically!

# Nonperiodic linear growth



The difference sequence  $\ell(n+1) - \ell(n)$  is

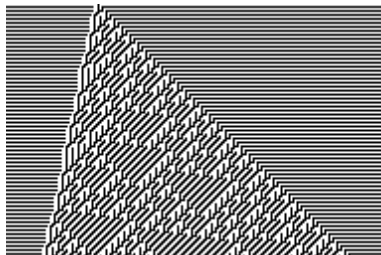
1, 2, 2, 1, 1, 2, 2, 1, 2, 2, 1, 1, 1, 2, 2, 1, 2, 2, 1, 1, 2, 2, 1, 2, 2, 1, 1, 1, 1, \dots

It can be obtained by dropping the first two letters of the fixed point

$$\varphi^\omega(2) = 2212211221221112212211221221111 \dots$$

of the morphism  $\varphi(1) = 1, \varphi(2) = 221$ .

# More nonperiodic linear growth



The difference sequence is  $(3\bar{1})^1(30)^2\psi(\varphi^\omega(A))$ , where

$$\varphi(A) = AC, \quad \varphi(B) = AD, \quad \varphi(C) = BA, \quad \varphi(D) = BB,$$

$$\psi(A) = (3\bar{1})^3(30)^2(3\bar{1})^3(30)^2(3\bar{1})^1(30)^2$$

$$\psi(B) = (3\bar{1})^3(30)^2(3\bar{1})^5(30)^2$$

$$\psi(C) = (3\bar{1})^5(30)^2(3\bar{1})^3(30)^2(3\bar{1})^1(30)^2$$

$$\psi(D) = (3\bar{1})^5(30)^2(3\bar{1})^5(30)^2.$$

## Definition

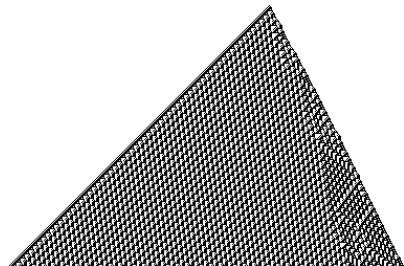
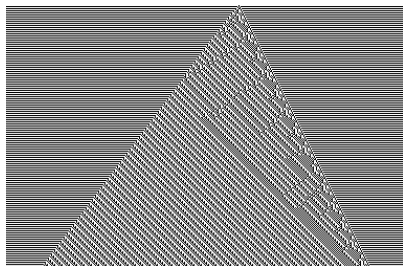
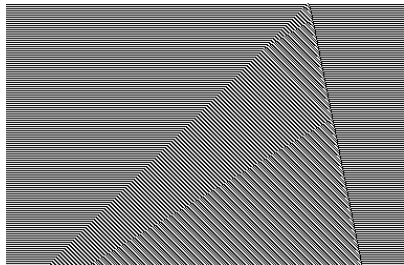
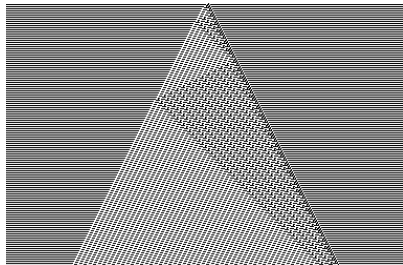
The **boundary word** is the infinite word  $(\ell(n+1) - \ell(n))_{n \geq 0}$ .

The boundary word is not necessarily a word on a finite subset of  $\mathbb{Z}$ .  
But often it is.

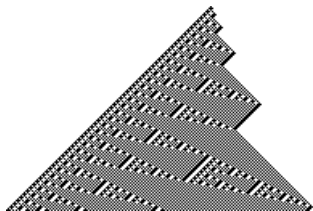
A classification of the 25088 automata is as follows.

- 24287 have eventually periodic boundary words.
- 81 have morphic boundary words that are not eventually periodic.
- 720 have boundaries that do not appear to be morphic.

# Oscillating particles



# Linear growth without a growth rate



Growth is linear:  $\ell(n) \in \Theta(n)$ .

But  $\lim_{n \rightarrow \infty} \frac{\ell(n)}{n}$  does not exist.

- $\liminf \ell(n)/n = 6/5$
- $\limsup \ell(n)/n = 3/2$

Boundary word is  $\psi(\varphi^\omega(A))$ , where

$$\varphi : A \rightarrow ABCB, B \rightarrow BB, C \rightarrow CC$$

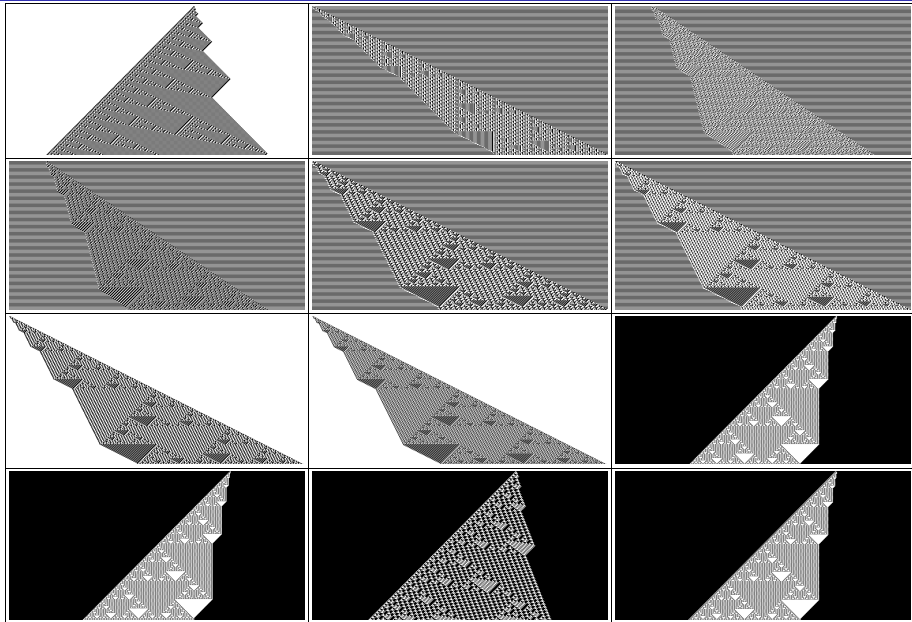
$$\psi : A \rightarrow \epsilon, B \rightarrow 2, C \rightarrow 0.$$

The fixed point of  $\varphi$  is

$$\varphi^\omega(A) = ABCBBBCCBBBBBBCCCCBBBBBBBBBBBBBB \dots$$

The frequencies of  $B$  and  $C$  don't exist!

# Automata with the same morphism $\varphi$

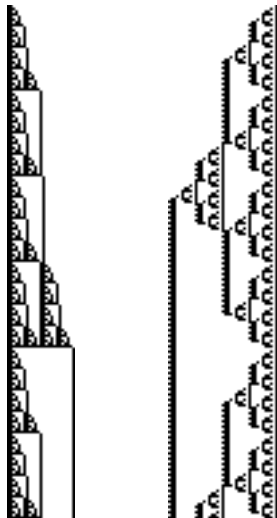




# Square-root growth

Some automata don't grow linearly:

- Rule 106 depending on 3 cells, begun from  $\dots \square\square\square \blacksquare \square\square\square \dots$ .
- Rule 39780 depending on 4 cells.



The boundary word for rule 106 is

$$\psi(\varphi^\omega(A)) = 11010011000000010000000011010011 \dots,$$

where

$$\varphi : A \rightarrow ABCD, B \rightarrow CCAB, C \rightarrow CCCC, D \rightarrow CCCD$$

$$\psi : A \rightarrow 1, B \rightarrow 1, C \rightarrow 0, D \rightarrow 1$$

Square-root growth rate can be derived from  $\varphi$ .

$\varphi$  is 4-uniform.



# Row lengths in rule 106

The length  $\ell(n)$  is 2-regular:

$$\ell(4n + 1) = 1/2\ell(4n) + 1/2\ell(4n + 2)$$

$$\ell(8n + 2) = -2\ell(2n) + \ell(8n) + 2\ell(2n + 1)$$

$$\ell(8n + 3) = -2\ell(2n) + \ell(8n) + 2\ell(2n + 1)$$

$$\ell(8n + 4) = -3\ell(2n) + \ell(8n) + 3\ell(2n + 1)$$

$$\ell(8n + 6) = -3\ell(2n) + \ell(8n) + 3\ell(2n + 1)$$

$$\ell(8n + 7) = -4\ell(2n) + \ell(8n) + 4\ell(2n + 1)$$

$$\ell(16n + 0) = -2\ell(n) + 3\ell(4n) + \ell(4n + 2) - \ell(4n + 3)$$

$$\ell(16n + 8) = -2\ell(n) + 1/2\ell(4n) + 7/2\ell(4n + 2) - \ell(4n + 3)$$



# $d = 4$ rule 39780

Rule 39780 also grows like  $\sqrt{n}$ .

Its boundary word is  $\psi(\varphi^\omega(A))$ , where

$$\varphi : A \rightarrow ABC, B \rightarrow DAB,$$

$$C \rightarrow CECE, D \rightarrow CECD, E \rightarrow CECE$$

$$\psi : A \rightarrow 2, B \rightarrow 2, C \rightarrow 1, D \rightarrow 0, E \rightarrow -1$$

$\varphi$  is not uniform.

$\ell(n)$  is evidently not 2-regular.



## Moral of the story

Properties of the automaton are reflected in the boundary word.

Vague principal:

Sufficiently simple boundaries are characterized by morphic words.